

Concurrent Algorithms

Exercise 10

(Given as an exam problem in 2016-2017)

Problem 1. Consider the following implementation of an obstruction-free consensus object from atomic multi-valued MRMW shared registers in a system of n processes. A process's id is known to itself as i .

Using: an array of atomic multi-valued MRMW shared registers $T[1, 2, \dots, n]$,
initialized to 0;

Using: an array of atomic multi-valued MRMW shared registers $V[1, 2, \dots, n]$,
initialized to $(\perp, 0)$;

```
propose( $v$ ) {
   $ts := i$ ;
  while (true) do{
     $T[i].write(ts)$ ;

     $maxts := 0$ ;
     $val := \perp$ ;
    for  $j = 1$  to  $n$  do
       $(t, vt) := V[j].read()$ ;
      if  $maxts < t$  then
         $maxts := t$ ;
         $val := vt$ ;

    if  $val = \perp$  then  $val := v$ ;

     $maxts := 0$ ;
    for  $j = 1$  to  $n$  do
       $t := T[j].read()$ ;
      if  $maxts < t$  then  $maxts := t$ ;

    if  $ts = maxts$  then
       $V[i].write(val, ts)$ ;
      return( $val$ );
     $ts := ts + n$ ;
  }
}
```

Recall that obstruction-free consensus ensures the property of *obstruction-freedom* instead of *wait-freedom*.

Your tasks:

1. Explain what is obstruction-freedom and what is the difference between obstruction-freedom and wait-freedom.
2. Answer whether the implementation satisfies obstruction-freedom. Justify your answer.
3. Is the algorithm correct? If the algorithm is correct prove its correctness. Otherwise, provide an execution that shows the algorithm is incorrect.